

# OpenMP Tutorial



Dirk Schmidl  
IT Center, RWTH Aachen University  
Member of the HPC Group  
[schmidl@itc.rwth-aachen.de](mailto:schmidl@itc.rwth-aachen.de)



Christian Terboven  
IT Center, RWTH Aachen University  
Head of the HPC Group  
[terboven@itc.rwth-aachen.de](mailto:terboven@itc.rwth-aachen.de)

# OpenMP Overview

***IWOMP Tutorial: October 5th, 2016***

**Christian Terboven**



# OpenMP Overview - Topics



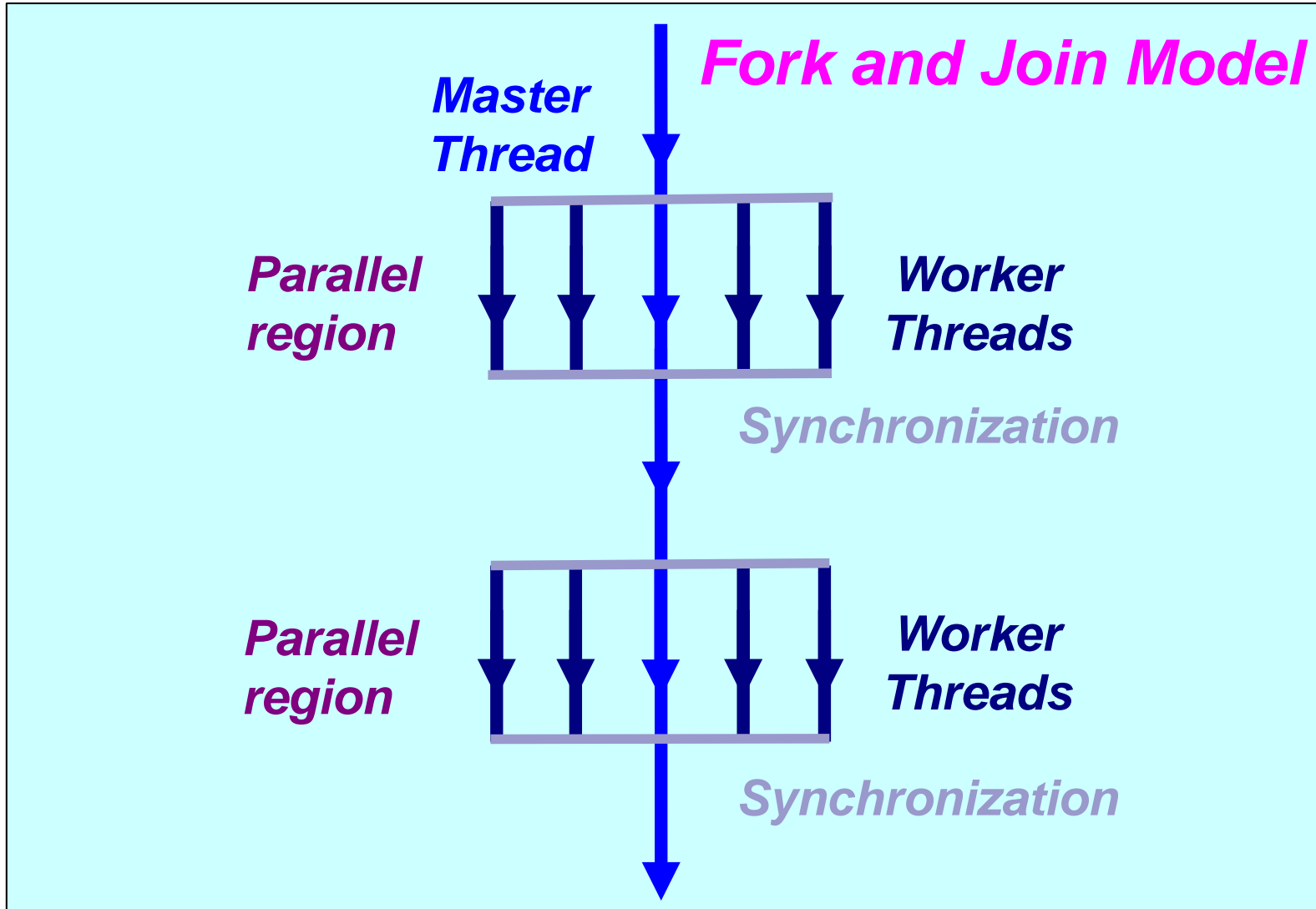
- Core Concepts
- Synchronization

*A parallel region is a block of code executed by all threads in the team*

```
#pragma omp parallel [clause[[,] clause] ...]
{
    "this code is executed in parallel"
} // End of parallel section (note: implied barrier)
```

```
!$omp parallel [clause[[,] clause] ...]
    "this code is executed in parallel"
!$omp end parallel (note: implied barrier)
```

# The OpenMP Execution Model



# The Worksharing Constructs

```
#pragma omp for
{
    ....
}
```

```
!$OMP DO
    ....
!$OMP END DO
```

```
#pragma omp sections
{
    ....
}
```

```
!$OMP SECTIONS
    ....
!$OMP END SECTIONS
```

```
#pragma omp single
{
    ....
}
```

```
!$OMP SINGLE
    ....
!$OMP END SINGLE
```

- ✓ *The work is distributed over the threads*
- ✓ *Must be enclosed in a parallel region*
- ✓ *Must be encountered by all threads in the team, or none at all*
- ✓ *No implied barrier on entry*
- ✓ *Implied barrier on exit (unless the nowait clause is specified)*
- ✓ *A work-sharing construct does not launch any new threads*

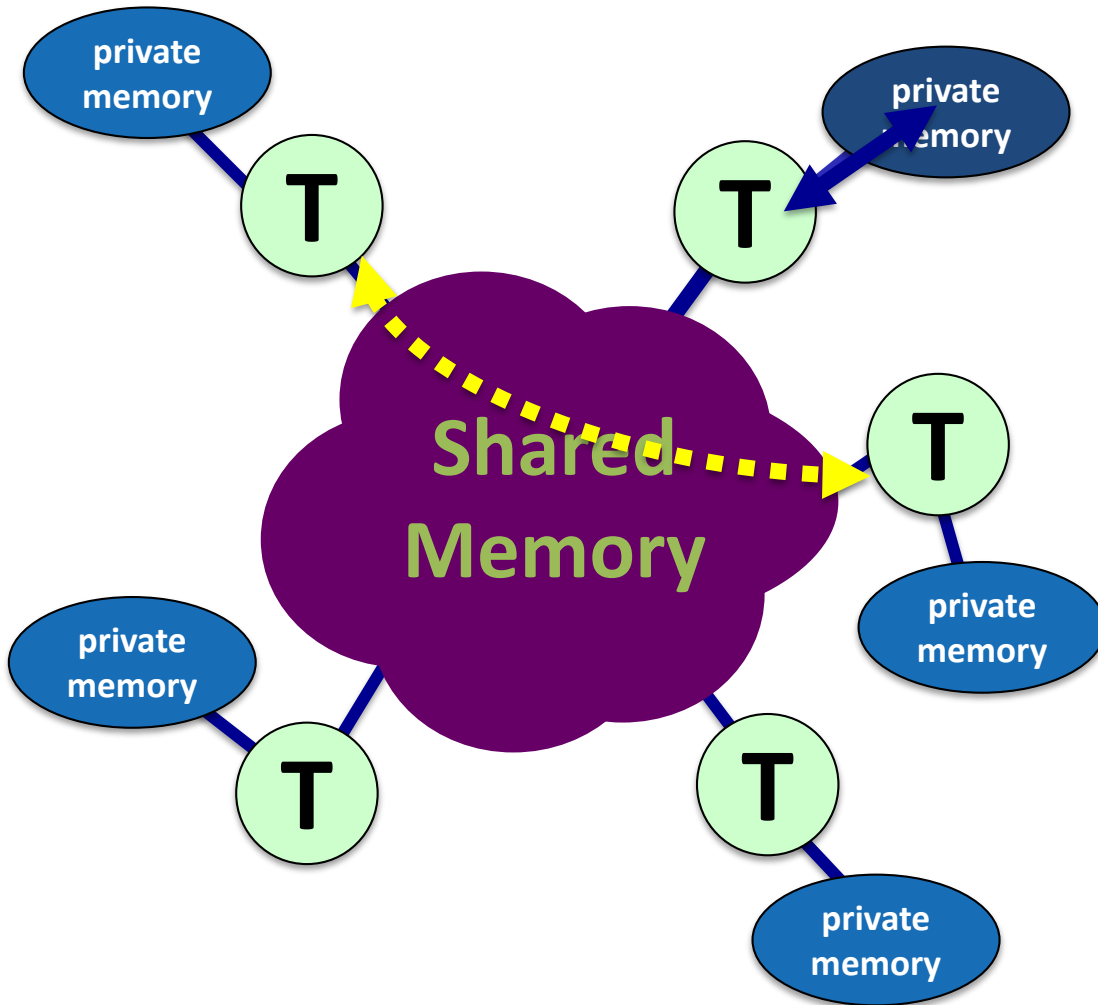
# The Single Directive

*Only one thread in the team executes the code enclosed*

```
#pragma omp single [private][firstprivate] \  
                  [copyprivate][nowait]  
{  
    <code-block>  
}
```

```
!$omp single [private][firstprivate]  
    <code-block>  
!$omp end single [copyprivate][nowait]
```

# The OpenMP Memory Model



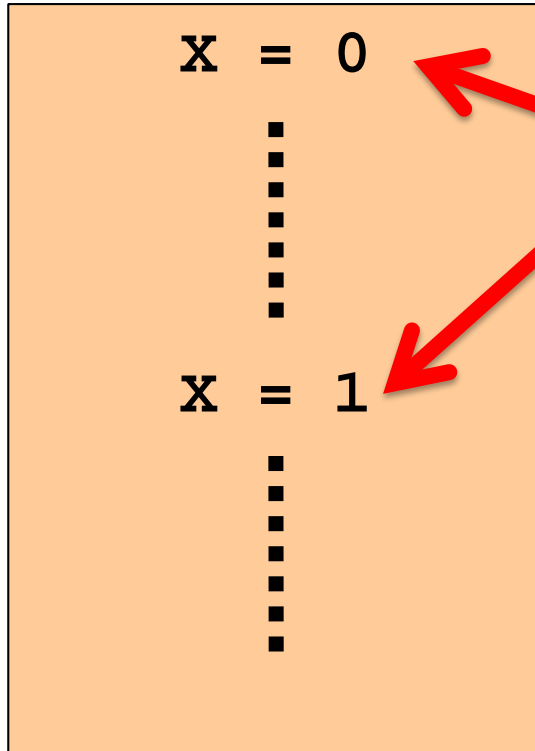
- ◆ *All threads have access to the same, globally shared memory*
- ◆ *Data in private memory is only accessible by the thread owning this memory*
- ◆ *No other thread sees the change(s) in private memory*
- ◆ *Data transfer is through shared memory and is 100% transparent to the application*



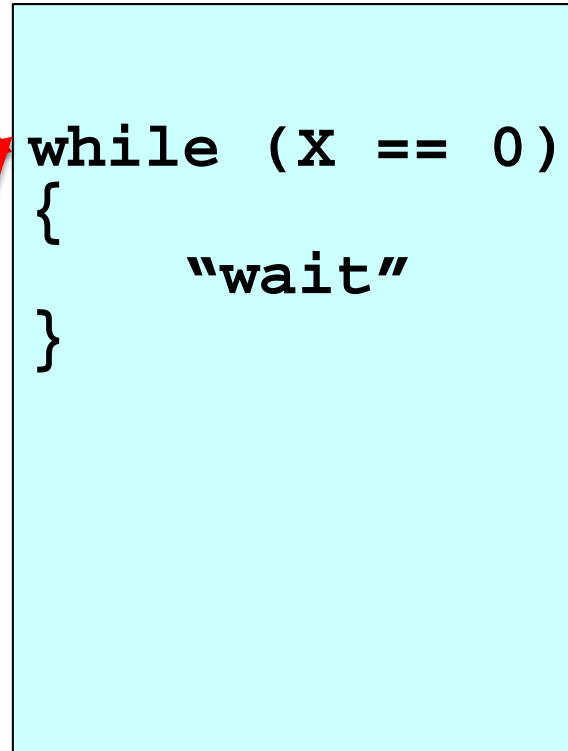
- Need to get this right
  - Part of the learning curve
- Private data is undefined on entry and exit
  - Can use `firstprivate` and `lastprivate` to address this
- Each thread has its own temporary view on the data
  - Applicable to shared data only
  - Means different threads may temporarily not see the same value for the same variable ...
  - Let me explain

# The Flush Directive

## Thread A



## Thread B



*If shared variable X is kept within a register, the modification may not be made visible to the other thread(s)*

# About The Flush

- Strongly recommended: do **not** use this directive
  - ... unless really necessary. Really 😊.
  - Could give very subtle interactions with compilers
  - If you insist on still doing so, be prepared to face the OpenMP language lawyers
  
- Implied on many constructs
  - A good thing
  - This is your safety net

# The OpenMP Barrier

- Several constructs have an implied barrier
  - This is another safety net (has implied flush by the way)
- In some cases, the implied barrier can be left out through the “nowait” clause
- This can help fine tuning the application
  - But you’d better know what you’re doing
- The explicit barrier comes in quite handy then

```
#pragma omp barrier
```

```
!$omp barrier
```

# The Nowait Clause

- To minimize synchronization, some directives support the optional nowait clause
  - If present, threads do not synchronize/wait at the end of that particular construct
- In C, it is one of the clauses on the pragma
- In Fortran, it is appended at the closing part of the construct

```
#pragma omp for nowait  
{  
    :  
}
```

```
!$omp do  
    :  
    :  
!$omp end do nowait
```